

SUPERQUBITS

Leron Borsten
Imperial College, London

in collaboration with:

K. Brádler, D. Dahanayake, M. J. Duff, H. Ebrahim, and W. Rubens
Phys. Rev. A 80, 032326 (2009)
arXiv:1206.6934

24 June - 3 July 2013
INTERNATIONAL SCHOOL OF SUBNUCLEAR PHYSICS
CCSEM, ERICE

Introduction

Superqubits

Super entanglement classification

Introduction

Entanglement and information

- In providing nonlocal resources, quantum mechanics distinguishes itself from the classical world.
- Perhaps the best known example of a nonlocal resource is the EPR-pair, introduced by Einstein, Podolsky and Rosen [1935]

$$\frac{1}{\sqrt{2}} (|01\rangle + |10\rangle)$$

Introduction

Entanglement and information

- In providing nonlocal resources, quantum mechanics distinguishes itself from the classical world.
- Perhaps the best known example of a nonlocal resource is the EPR-pair, introduced by Einstein, Podolsky and Rosen [1935]

$$\frac{1}{\sqrt{2}} (|01\rangle + |10\rangle)$$

- Bell famously showed that this quantum state can be used to violate an inequality required to be satisfied by *any* local hidden variable model [Bell:1964].

Introduction

Entanglement and information

- In providing nonlocal resources, quantum mechanics distinguishes itself from the classical world.
- Perhaps the best known example of a nonlocal resource is the EPR-pair, introduced by Einstein, Podolsky and Rosen [1935]

$$\frac{1}{\sqrt{2}} (|01\rangle + |10\rangle)$$

- Bell famously showed that this quantum state can be used to violate an inequality required to be satisfied by *any* local hidden variable model [Bell:1964].
- This fundamental shift in our understanding of reality has more recently precipitated a generalisation of *information theory*: *qubits*

$$0 \text{ or } 1 \quad \rightarrow \quad a_A |A\rangle = a_0 |0\rangle + a_1 |1\rangle \in \mathbb{C}^2$$

$$0010 \dots 0101 \quad \rightarrow \quad a_{A_1 A_2 \dots A_n} |A_1 A_2 \dots A_n\rangle \in \mathbb{C}^2 \otimes \mathbb{C}^2 \dots \otimes \mathbb{C}^2$$

- Entanglement classification is an essential open problem

Introduction

Supersymmetric quantum information

- Here we propose a supersymmetric generalization of the qubit, the *superqubit*

$$|\Psi\rangle = |A\rangle a_A + |\bullet\rangle a_\bullet$$

- A single superqubit forms a 3-dimensional representation of $OSp(1|2)$ consisting of two commuting “bosonic” components and one anticommuting “fermionic” component.

Introduction

Supersymmetric quantum information

- Here we propose a supersymmetric generalization of the qubit, the *superqubit*

$$|\Psi\rangle = |A\rangle a_A + |\bullet\rangle a_\bullet$$

- A single superqubit forms a 3-dimensional representation of $OSp(1|2)$ consisting of two commuting “bosonic” components and one anticommuting “fermionic” component.

Superqubit entanglement

- Super-Bell and super-GHZ states are characterized, respectively, by nonvanishing superdeterminant (distinct from the Berezinian) and superhyperdeterminant

Entanglement: Non-local games

The 3-player game

[Greenberger, Horne, Zeilinger: 1989; Mermin: 1990; Watrous et al: 2004]

- *Players* (Alice, Bob, Charlie. . .): act cooperatively in order to win.
- *Referee*: coordinates the game, question set \mathcal{Q} and answer set \mathcal{A}

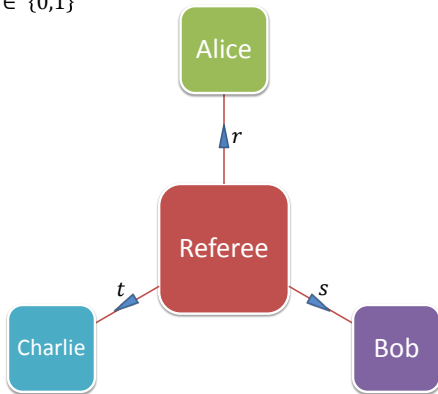
Entanglement: Non-local games

The 3-player game

[Greenberger, Horne, Zeilinger: 1989; Mermin: 1990; Watrous et al: 2004]

- *Players* (Alice, Bob, Charlie...): act cooperatively in order to win.
- *Referee*: coordinates the game, question set \mathcal{Q} and answer set \mathcal{A}

$$r, s, t \in \{0,1\}$$



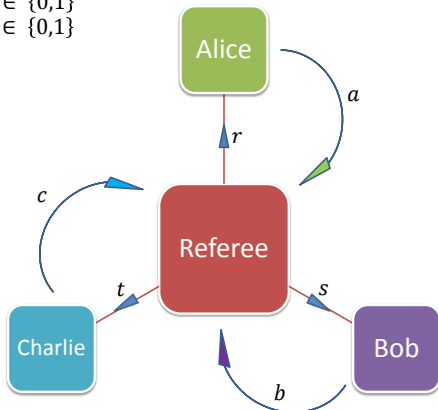
Entanglement: Non-local games

The 3-player game

[Greenberger, Horne, Zeilinger: 1989; Mermin: 1990; Watrous et al: 2004]

- *Players* (Alice, Bob, Charlie...): act cooperatively in order to win.
- *Referee*: coordinates the game, question set \mathcal{Q} and answer set \mathcal{A}

$$r, s, t \in \{0,1\}$$
$$a, b, c \in \{0,1\}$$



Entanglement: Non-local games

The 3-player game: rules for winning

- The referee ensures that $rst \in \{000, 110, 101, 011\}$ and the players are aware of this.
- The players win if $r \vee s \vee t = a \oplus b \oplus c$, where \vee and \oplus respectively denote disjunction and addition mod 2:

rs	$a \oplus b \oplus c$
000	0
011	1
101	1
110	1

The 3-player game: Deterministic strategy

- What is the best possible classical *deterministic* strategy?
- A deterministic strategy amounts to specifying three functions, one for each player, from the question set \mathcal{Q} to the answer set \mathcal{A}

$$a : \mathcal{Q} \rightarrow \mathcal{A}; \quad r \mapsto a(r),$$

$$b : \mathcal{Q} \rightarrow \mathcal{A}; \quad s \mapsto b(s),$$

$$c : \mathcal{Q} \rightarrow \mathcal{A}; \quad t \mapsto c(t),$$

The 3-player game: Deterministic strategy

- What is the best possible classical *deterministic* strategy?
- A deterministic strategy amounts to specifying three functions, one for each player, from the question set \mathcal{Q} to the answer set \mathcal{A}

$$a : \mathcal{Q} \rightarrow \mathcal{A}; \quad r \mapsto a(r),$$

$$b : \mathcal{Q} \rightarrow \mathcal{A}; \quad s \mapsto b(s),$$

$$c : \mathcal{Q} \rightarrow \mathcal{A}; \quad t \mapsto c(t),$$

The condition that the players win may then be written as,

$$a(0) \oplus b(0) \oplus c(0) = 0,$$

$$a(1) \oplus b(1) \oplus c(0) = 1,$$

$$a(1) \oplus b(0) \oplus c(1) = 1,$$

$$a(0) \oplus b(1) \oplus c(1) = 1.$$

- Add mod 2 \rightarrow contradiction \Rightarrow best one can do is win 75% of the time

The 3-player game: Deterministic strategy

- What is the best possible classical *deterministic* strategy?
- A deterministic strategy amounts to specifying three functions, one for each player, from the question set \mathcal{Q} to the answer set \mathcal{A}

$$a : \mathcal{Q} \rightarrow \mathcal{A}; \quad r \mapsto a(r),$$

$$b : \mathcal{Q} \rightarrow \mathcal{A}; \quad s \mapsto b(s),$$

$$c : \mathcal{Q} \rightarrow \mathcal{A}; \quad t \mapsto c(t),$$

The condition that the players win may then be written as,

$$a(0) \oplus b(0) \oplus c(0) = 0,$$

$$a(1) \oplus b(1) \oplus c(0) = 1,$$

$$a(1) \oplus b(0) \oplus c(1) = 1,$$

$$a(0) \oplus b(1) \oplus c(1) = 1.$$

- Add mod 2 \rightarrow contradiction \Rightarrow best one can do is win 75% of the time

In this language the contradiction with local realism exposed by Mermin translates into the existence of a local strategy utilising the GHZ state that wins the game with $p_{win} > 0.75$.

The 3-player game: Quantum strategy

Alice, Bob and Charlie share an entangled GHZ state:

$$|\Psi\rangle = \frac{1}{\sqrt{2}} (|000\rangle - |011\rangle - |101\rangle - |110\rangle)$$

The 3-player game: Quantum strategy

Alice, Bob and Charlie share an entangled GHZ state:

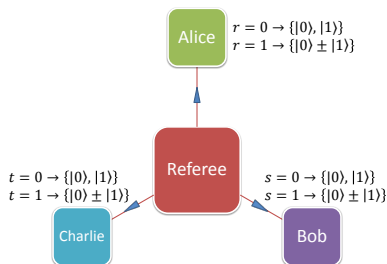
$$|\Psi\rangle = \frac{1}{\sqrt{2}} (|000\rangle - |011\rangle - |101\rangle - |110\rangle)$$

- If $r, s, t = 0$ then Alice, Bob, Charlie measure in the computational basis

$$\{|0\rangle, |1\rangle\}$$

- If $r, s, t = 1$ then Alice, Bob, Charlie measure in the Hadamard basis

$$\left\{ \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle), \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) \right\}$$



The 3-player game: Quantum strategy

Winning probability

1. $rst = 000$: All measure in the computational basis

$$|\Psi\rangle = \frac{1}{\sqrt{2}} (|000\rangle - |011\rangle - |101\rangle - |110\rangle)$$

$\Rightarrow a \oplus b \oplus c = 0$. Always win.

The 3-player game: Quantum strategy

Winning probability

1. $rst = 000$: All measure in the computational basis

$$|\Psi\rangle = \frac{1}{\sqrt{2}} (|000\rangle - |011\rangle - |101\rangle - |110\rangle)$$

$\Rightarrow a \oplus b \oplus c = 0$. Always win.

2. $rst = \{011, 101, 110\}$: Two measure in the Hadamard basis

$$\mathbb{1} \otimes H \otimes H |\Psi\rangle = \frac{1}{\sqrt{2}} (|001\rangle + |010\rangle - |100\rangle + |111\rangle)$$

$\Rightarrow a \oplus b \oplus c = 1$. Always win.

The 3-player game: Quantum strategy

Winning probability

1. $rst = 000$: All measure in the computational basis

$$|\Psi\rangle = \frac{1}{\sqrt{2}} (|000\rangle - |011\rangle - |101\rangle - |110\rangle)$$

$\Rightarrow a \oplus b \oplus c = 0$. Always win.

2. $rst = \{011, 101, 110\}$: Two measure in the Hadamard basis

$$\mathbb{1} \otimes H \otimes H |\Psi\rangle = \frac{1}{\sqrt{2}} (|001\rangle + |010\rangle - |100\rangle + |111\rangle)$$

$\Rightarrow a \oplus b \oplus c = 1$. Always win.

Conclusion

With an entangled resource we can win the game with certainty:

At the tables quantum players beat their poor classical cousins!

Different kinds of entanglement

Comparing entanglement properties

- Using the totally entangled state GHZ state

$$|\Psi\rangle_{GHZ} = \frac{1}{\sqrt{2}} (|000\rangle - |011\rangle - |101\rangle - |110\rangle)$$

With a single run we can rule out local realist hidden model theories

- However, playing the same game this totally entangled state

$$|\Psi\rangle_W = \frac{1}{\sqrt{2}} (|001\rangle + |010\rangle + |100\rangle)$$

we only win with $p_{win}(W) = 7/8$. Cannot win with certainty [LB:2013]

Different kinds of entanglement

Comparing entanglement properties

- Using the totally entangled state GHZ state

$$|\Psi\rangle_{GHZ} = \frac{1}{\sqrt{2}} (|000\rangle - |011\rangle - |101\rangle - |110\rangle)$$

With a single run we can rule out local realist hidden model theories

- However, playing the same game this totally entangled state

$$|\Psi\rangle_W = \frac{1}{\sqrt{2}} (|001\rangle + |010\rangle + |100\rangle)$$

we only win with $p_{win}(W) = 7/8$. Cannot win with certainty [LB:2013]

- How are we to mathematically distinguish the properties of W and GHZ?

Measuring and classifying entanglement

Entanglement classes

- Two n -qubit states have the same entanglement iff they are related by the equivalence group of *Stochastic Local Operations and Classical Communication* (SLOCC)

$$[\mathrm{SL}(2, \mathbb{C})]^{\otimes n}$$

- “gauge” group of n -qubit entanglement
- The space of entanglement classes:

$$\frac{\mathbb{C}^2 \otimes \mathbb{C}^2 \otimes \dots \otimes \mathbb{C}^2}{\mathrm{SL}_1(2, \mathbb{C}) \times \mathrm{SL}_2(2, \mathbb{C}) \times \dots \times \mathrm{SL}_n(2, \mathbb{C})}$$

[Bennett et al:1999, Dur et al:2000]

2-qubit entanglement classification

$$|\psi\rangle = a_{AB}|AB\rangle$$

2-qubit entanglement classification

$$|\psi\rangle = a_{AB}|AB\rangle$$

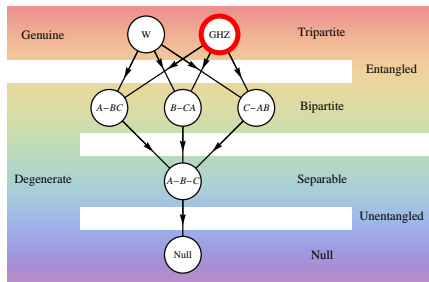
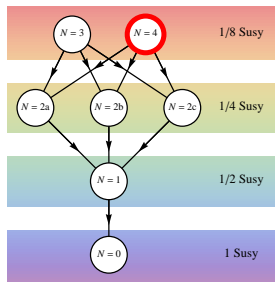
Just two entanglement classes separated by the $SL_A(2, \mathbb{C}) \times SL_A(2, \mathbb{C})$ -invariant

$$\det a_{AB}$$

1. Separable A - B : $\det a_{AB} = 0$
2. Entangled EPR: $\det a_{AB} \neq 0$

3-qubit entanglement classification

$$|\psi\rangle = a_{ABC}|ABC\rangle$$



- There are 6 entanglement classes [Dur-Vidal:2000]
- SLOCC \rightarrow Freudenthal Triple Systems and “Groups of type E_7 ”: HA/AH row/column the magic square, cf. MJ Duff’s lectures [Ferrara-Gunaydin:1997; LB-Dahanayake-Duff-Ebrahim-Rubens:2008, 2009]

3-qubit entanglement classification

Six entanglement classes are separated by $[\mathrm{SL}(2, \mathbb{C})]^3$ -covariants
[LB-Dahanayake-Duff-Ebrahim-Rubens:2009]:

1. Separable A - B - C : $\gamma^A = \gamma^B = \gamma^C = 0$

$$(\gamma^A)_{A_1 A_2} := a_{A_1}{}^{BC} a_{A_2 BC}, \quad (\gamma^B)_{B_1 B_2} := a_{B_1}{}^A C a_{AB_2 C}, \quad (\gamma^C)_{C_1 C_2} := a^{AB}{}_{C_1} a_{ABC_2},$$

3-qubit entanglement classification

Six entanglement classes are separated by $[\mathrm{SL}(2, \mathbb{C})]^3$ -covariants
 [LB-Dahanayake-Duff-Ebrahim-Rubens:2009]:

1. Separable A - B - C : $\gamma^A = \gamma^B = \gamma^C = 0$

$$(\gamma^A)_{A_1 A_2} := a_{A_1}{}^{BC} a_{A_2 BC}, \quad (\gamma^B)_{B_1 B_2} := a_{B_1}{}^A C a_{AB_2 C}, \quad (\gamma^C)_{C_1 C_2} := a^{AB}{}_{C_1} a_{ABC_2},$$

2. Three biseparable A - EPR : $\gamma^A \neq 0$ and $\gamma^B = \gamma^C = 0$

3-qubit entanglement classification

Six entanglement classes are separated by $[\mathrm{SL}(2, \mathbb{C})]^3$ -covariants
 [LB-Dahanayake-Duff-Ebrahim-Rubens:2009]:

1. Separable A - B - C : $\gamma^A = \gamma^B = \gamma^C = 0$

$$(\gamma^A)_{A_1 A_2} := a_{A_1}{}^{BC} a_{A_2 BC}, \quad (\gamma^B)_{B_1 B_2} := a^A{}_{B_1}{}^C a_{AB_2 C}, \quad (\gamma^C)_{C_1 C_2} := a^{AB}{}_{C_1} a_{ABC_2},$$

2. Three biseparable A - EPR : $\gamma^A \neq 0$ and $\gamma^B = \gamma^C = 0$

3. Totally entangled W states: $T_{ABC} \neq 0$ and $\mathrm{Det} a_{ABC} = 0$

$$T_{ABC} := (\gamma^A)_{AA'} a^{A'}{}_{BC}$$

and $\mathrm{Det} a$ is Cayley's **Hyperdeterminant** [Cayley:1845]

$$\mathrm{Det} a_{ABC} := -\det \gamma^A = -\det \gamma^B = -\det \gamma^C$$

3-qubit entanglement classification

Six entanglement classes are separated by $[\mathrm{SL}(2, \mathbb{C})]^3$ -covariants
 [LB-Dahanayake-Duff-Ebrahim-Rubens:2009]:

1. Separable A - B - C : $\gamma^A = \gamma^B = \gamma^C = 0$

$$(\gamma^A)_{A_1 A_2} := a_{A_1}{}^{BC} a_{A_2 BC}, \quad (\gamma^B)_{B_1 B_2} := a^A{}_{B_1}{}^C a_{AB_2 C}, \quad (\gamma^C)_{C_1 C_2} := a^{AB}{}_{C_1} a_{ABC_2},$$

2. Three biseparable A - EPR : $\gamma^A \neq 0$ and $\gamma^B = \gamma^C = 0$

3. Totally entangled W states: $T_{ABC} \neq 0$ and $\mathrm{Det} a_{ABC} = 0$

$$T_{ABC} := (\gamma^A)_{AA'} a^{A'}{}_{BC}$$

and $\mathrm{Det} a$ is Cayley's **Hyperdeterminant** [Cayley:1845]

$$\mathrm{Det} a_{ABC} := -\det \gamma^A = -\det \gamma^B = -\det \gamma^C$$

4. Totally entangled GHZ states: $\mathrm{Det} a_{ABC} \neq 0$

3-qubit entanglement classification

Six entanglement classes are separated by $[\mathrm{SL}(2, \mathbb{C})]^3$ -covariants
 [LB-Dahanayake-Duff-Ebrahim-Rubens:2009]:

1. Separable A - B - C : $\gamma^A = \gamma^B = \gamma^C = 0$

$$(\gamma^A)_{A_1 A_2} := a_{A_1}{}^{BC} a_{A_2 BC}, \quad (\gamma^B)_{B_1 B_2} := a^A{}_{B_1}{}^C a_{AB_2 C}, \quad (\gamma^C)_{C_1 C_2} := a^{AB}{}_{C_1} a_{ABC_2},$$

2. Three biseparable A - EPR : $\gamma^A \neq 0$ and $\gamma^B = \gamma^C = 0$

3. Totally entangled W states: $T_{ABC} \neq 0$ and $\mathrm{Det} a_{ABC} = 0$

$$T_{ABC} := (\gamma^A)_{AA'} a^{A'}{}_{BC}$$

and $\mathrm{Det} a$ is Cayley's **Hyperdeterminant** [Cayley:1845]

$$\mathrm{Det} a_{ABC} := -\det \gamma^A = -\det \gamma^B = -\det \gamma^C$$

4. Totally entangled GHZ states: $\mathrm{Det} a_{ABC} \neq 0$

$\mathrm{Det} a = 0 \Rightarrow$ no strategy that wins the 3-player game with certainty [LB:2013]

Introduction

Superqubits

Super entanglement classification

The superqubit

The superqubit [LB-Dahanayake-Duff-Rubens: 2009]:

$$\begin{aligned} |\Psi\rangle &= |X\rangle a_X \\ &= |A\rangle a_A + |\bullet\rangle a_\bullet \end{aligned}$$

where a_A ($A = 0, 1$) and a_\bullet are commuting/anticommuting supernumbers

The superqubit

The superqubit [LB-Dahanayake-Duff-Rubens: 2009]:

$$\begin{aligned} |\Psi\rangle &= |X\rangle a_X \\ &= |A\rangle a_A + |\bullet\rangle a_\bullet \end{aligned}$$

where a_A ($A = 0, 1$) and a_\bullet are commuting/anticommuting supernumbers

- Even element of a *super Hilbert space* [DeWitt:1984, Rogers:1980, Rudolph:2000]:

$$\dagger : \mathcal{H} \rightarrow \mathcal{H}^\dagger; \quad |\psi\rangle \mapsto (|\psi\rangle)^\dagger := \langle\psi|.$$

1. \dagger sends pure bosonic (fermionic) supervectors in \mathcal{H} into bosonic (fermionic) supervectors in \mathcal{H}^\dagger .
2. \dagger is linear

$$(|\psi\rangle + |\phi\rangle)^\dagger = \langle\psi| + \langle\phi|.$$

3. For pure even/odd α and $|\psi\rangle$

$$(|\psi\rangle\alpha)^\dagger = (-)^{\alpha\psi} \alpha^\# \langle\psi|, \quad (\alpha\langle\psi|)^\dagger = (-)^{\psi+\alpha\psi} |\psi\rangle\alpha^\#,$$

where $\#$ is the superconjugate.

Superqubit norm

The even Grassmann valued norm

$$\langle \Psi | \Psi \rangle = \delta^{A_1 A_2} a_{A_1}^\# a_{A_2} - a_{\bullet}^\# a_{\bullet},$$

is invariant under the *supergroup* of local unitaries

$$\mathrm{SU}(2) \longrightarrow \mathrm{UOSp}(2|1)$$

[Berezin-Tolstoy:1981]

- A normalized superqubit $|\Psi\rangle$ may be regarded as an element of the projective space [Landi:1999]

$$S^{2|2} = \mathrm{UOSp}(1|2) / \mathrm{U}(0|1)$$

known as the *supersphere*

$$|\Psi\rangle = Z(\eta, \alpha, \beta)|0\rangle, \quad \text{where } Z \in \mathrm{UOSp}(2|1)$$

$$Z(\eta, \alpha, \beta) = S(\eta)U(\alpha, \beta) = \begin{pmatrix} 1 + \frac{1}{4}\eta\eta^\# & -\frac{\eta}{2} & \frac{\eta^\#}{2} \\ -\frac{\eta^\#}{2} & 1 - \frac{1}{8}\eta\eta^\# & 0 \\ -\frac{\eta}{2} & 0 & 1 - \frac{1}{8}\eta\eta^\# \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 \\ 0 & \alpha & -\beta^\# \\ 0 & \beta & \alpha^\# \end{pmatrix}$$

Superqubits

2-superqubit

$$|\Psi\rangle = |AB\rangle a_{AB} + |A\bullet\rangle a_{A\bullet} + |\bullet B\rangle a_{\bullet B} + |\bullet\bullet\rangle a_{\bullet\bullet}$$

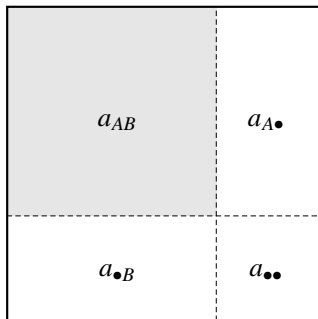


Figure : The 3×3 square supermatrix

Superqubits

3-superqubit

$$\begin{aligned}
 |\Psi\rangle = & |ABC\rangle a_{ABC} + |AB\bullet\rangle a_{AB\bullet} + |A\bullet C\rangle a_{A\bullet C} + |\bullet BC\rangle a_{\bullet BC} \\
 & + |A\bullet\bullet\rangle a_{A\bullet\bullet} + |\bullet B\bullet\rangle a_{\bullet B\bullet} + |\bullet\bullet C\rangle a_{\bullet\bullet C} \\
 & + |\bullet\bullet\bullet\rangle a_{\bullet\bullet\bullet}
 \end{aligned}$$

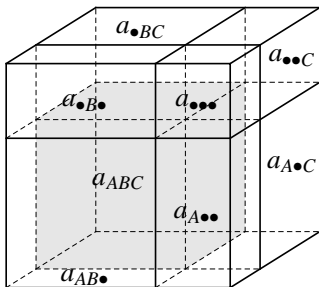


Figure : The $3 \times 3 \times 3$ cubic superhypermatrix

Introduction

Superqubits

Super entanglement classification

Super entanglement classification

The super SLOCC group

$$\boxed{\mathrm{SL}(2, \mathbb{C}) \longrightarrow \mathrm{OSp}(2|1)}$$

[Berezin-Tolstoy:1981]

Super entanglement classification

The super SLOCC group

$$\boxed{\mathrm{SL}(2, \mathbb{C}) \longrightarrow \mathrm{OSp}(2|1)}$$

[Berezin-Tolstoy:1981]

2 superqubits: the superdeterminant

$$\mathrm{SL}_A(2, \mathbb{C}) \times \mathrm{SL}_B(2, \mathbb{C}) \longrightarrow \mathrm{OSp}_A(2|1) \times \mathrm{OSp}_B(2|1)$$

$\mathrm{OSp}_A(2|1) \times \mathrm{OSp}_B(2|1)$ -invariant superdeterminant

$$\begin{aligned} \mathrm{sdet} a_{XY} &= \frac{1}{2} (a^{AB} a_{AB} - a^{A\bullet} a_{A\bullet} - a^{\bullet B} a_{\bullet B} - a^{\bullet\bullet} a_{\bullet\bullet}) \\ &= (a_{00} a_{11} - a_{01} a_{10} + a_{0\bullet} a_{1\bullet} + a_{\bullet 0} a_{\bullet 1}) - \frac{1}{2} a_{\bullet\bullet}^2, \end{aligned}$$

Super entanglement classification

The super SLOCC group

$$\boxed{\mathrm{SL}(2, \mathbb{C}) \longrightarrow \mathrm{OSp}(2|1)}$$

[Berezin-Tolstoy:1981]

2 superqubits: the superdeterminant

$$\mathrm{SL}_A(2, \mathbb{C}) \times \mathrm{SL}_B(2, \mathbb{C}) \longrightarrow \mathrm{OSp}_A(2|1) \times \mathrm{OSp}_B(2|1)$$

$\mathrm{OSp}_A(2|1) \times \mathrm{OSp}_B(2|1)$ -invariant superdeterminant

$$\begin{aligned} \mathrm{sdet} a_{XY} &= \frac{1}{2} (a^{AB} a_{AB} - a^{A\bullet} a_{A\bullet} - a^{\bullet B} a_{\bullet B} - a^{\bullet\bullet} a_{\bullet\bullet}) \\ &= (a_{00} a_{11} - a_{01} a_{10} + a_{0\bullet} a_{1\bullet} + a_{\bullet 0} a_{\bullet 1}) - \frac{1}{2} a_{\bullet\bullet}^2, \end{aligned}$$

Not the Berezinian, but quite natural:

$$\det a_{AB} = \frac{1}{2} \mathrm{tr}[(a\varepsilon)^t \varepsilon a] \longrightarrow \mathrm{sdet} a_{XY} = \frac{1}{2} \mathrm{str}[(aE)^{st} Ea]$$

3 superqubits: the superhyperdeterminant

Recall, for 3 qubits we had $[\mathrm{SL}(2, \mathbb{C})]^3$ -covariants

$$a_{ABC}$$

$$\gamma_{A_1 A_2}^A$$

$$\gamma_{B_1 B_2}^B$$

$$\gamma_{C_1 C_2}^C$$

$$T_{ABC}$$

$$\mathrm{Det} a_{ABC}$$

3 superqubits: the superhyperdeterminant

Need $[\mathrm{OSp}(2|1)]^3$ -supercovariants:

1. Quadratic $(\mathbf{5}, \mathbf{1}, \mathbf{1})$

$$\Gamma_{X_1 X_2}^A := \left(\begin{array}{c|c} \gamma_{A_1 A_2} & \gamma_{A_1 \bullet} \\ \hline \gamma_{\bullet A_2} & \gamma_{\bullet \bullet} \end{array} \right) = \left(\begin{array}{c|c} \gamma_{A_1 A_2} & \gamma_{A_1 \bullet} \\ \hline \gamma_{A_2 \bullet} & 0 \end{array} \right),$$

$$\gamma_{A_1 A_2} := a_{A_1}{}^{BC} a_{A_2 BC} - a_{A_1}{}^{B\bullet} a_{A_2 B\bullet} - a_{A_1}{}^{\bullet C} a_{A_2 \bullet C} - a_{A_1}{}^{\bullet\bullet} a_{A_2 \bullet\bullet}$$

$$\gamma_{A_1 \bullet} := a_{A_1}{}^{BC} a_{\bullet BC} + a_{A_1}{}^{B\bullet} a_{\bullet B\bullet} + a_{A_1}{}^{\bullet C} a_{\bullet\bullet C} - a_{A_1}{}^{\bullet\bullet} a_{\bullet\bullet\bullet}$$

2. Cubic triple product $(\mathbf{3}, \mathbf{3}, \mathbf{3})$

$$T_{XYZ} = \Gamma_{XX'}^A a^{X' YZ}.$$

3. Quartic superhyperdeterminant $(\mathbf{1}, \mathbf{1}, \mathbf{1})$

$$\mathrm{sDet} a_{XYZ} = -\mathrm{sdet} \Gamma^A = -\mathrm{sdet} \Gamma^B = -\mathrm{sdet} \Gamma^C$$

This is the unique quartic invariant first found in [\[Castellani-Grassi-Sommovigo:2010\]](#) in the context of $\mathcal{N} = 2$ sugra.

Summary

3-superqubit entanglement

- All 3-qubit entanglement covariants may supersymmetrized in a very nature manner:

$$\begin{array}{lll}
 a_{ABC} & \rightarrow & a_{XYZ} \\
 \gamma_A & \rightarrow & \Gamma_X \\
 \gamma_B & \rightarrow & \Gamma_Y \\
 \gamma_C & \rightarrow & \Gamma_Z \\
 T_{ABC} & \rightarrow & T_{XYZ} \\
 \text{Det } a_{ABC} & \rightarrow & \text{sDet } a_{XYZ}
 \end{array}$$

such that

$$\text{sDet } a_{XYZ} = -\text{sdet } \Gamma_X = -\text{sdet } \Gamma_Y = -\text{sdet } \Gamma_Z.$$

cf.

$$\text{Det } a_{ABC} = -\det \Gamma_A = -\det \Gamma_B = -\det \Gamma_C.$$

- Cayley's Hyperdeterminant may be supersymmetrized! The **Superhyperdeterminant** is invariant under $\text{OSp}(2|1) \times \text{OSp}(2|1) \times \text{OSp}(2|1)$

Conclusions

- Here we proposed a supersymmetric generalization of the qubit, the *superqubit*.
- Not HEP susy:
 - Not a representation of the super-Poincaré group
 - It is a non-trivial *consistent* extension of conventional quantum mechanics based on a super Hilbert space [DeWitt:1984, Rogers:1980, Rudolph:2000]

Conclusions

- Here we proposed a supersymmetric generalization of the qubit, the *superqubit*.
- Not HEP susy:
 - Not a representation of the super-Poincaré group
 - It is a non-trivial *consistent* extension of conventional quantum mechanics based on a super Hilbert space [DeWitt:1984, Rogers:1980, Rudolph:2000]
- Extend the n -qubit SLOCC equivalence group $[SL(2, \mathbb{C})]^n$ and the LOCC equivalence group $[SU(2)]^n$ to the supergroups $[OSp(1|2)]^n$ and $[UOSp(1|2)]^n$, respectively.

Conclusions

- Here we proposed a supersymmetric generalization of the qubit, the *superqubit*.
- Not HEP susy:
 - Not a representation of the super-Poincaré group
 - It is a non-trivial *consistent* extension of conventional quantum mechanics based on a super Hilbert space [DeWitt:1984, Rogers:1980, Rudolph:2000]
- Extend the n -qubit SLOCC equivalence group $[SL(2, \mathbb{C})]^n$ and the LOCC equivalence group $[SU(2)]^n$ to the supergroups $[OSp(1|2)]^n$ and $[UOSp(1|2)]^n$, respectively.
- For $n = 2$ and $n = 3$ we introduce the appropriate supersymmetric generalizations of the conventional entanglement measures.

Conclusions

- Here we proposed a supersymmetric generalization of the qubit, the *superqubit*.
- Not HEP susy:
 - Not a representation of the super-Poincaré group
 - It is a non-trivial *consistent* extension of conventional quantum mechanics based on a super Hilbert space [DeWitt:1984, Rogers:1980, Rudolph:2000]
- Extend the n -qubit SLOCC equivalence group $[SL(2, \mathbb{C})]^n$ and the LOCC equivalence group $[SU(2)]^n$ to the supergroups $[OSp(1|2)]^n$ and $[UOSp(1|2)]^n$, respectively.
- For $n = 2$ and $n = 3$ we introduce the appropriate supersymmetric generalizations of the conventional entanglement measures.
- In particular, super-Bell and super-GHZ states are characterized, respectively, by nonvanishing superdeterminant (distinct from the Berezinian) and superhyperdeterminant

Conclusions

- Here we proposed a supersymmetric generalization of the qubit, the *superqubit*.
- Not HEP susy:
 - Not a representation of the super-Poincaré group
 - It is a non-trivial *consistent* extension of conventional quantum mechanics based on a super Hilbert space [DeWitt:1984, Rogers:1980, Rudolph:2000]
- Extend the n -qubit SLOCC equivalence group $[SL(2, \mathbb{C})]^n$ and the LOCC equivalence group $[SU(2)]^n$ to the supergroups $[OSp(1|2)]^n$ and $[UOSp(1|2)]^n$, respectively.
- For $n = 2$ and $n = 3$ we introduce the appropriate supersymmetric generalizations of the conventional entanglement measures.
- In particular, super-Bell and super-GHZ states are characterized, respectively, by nonvanishing superdeterminant (distinct from the Berezinian) and superhyperdeterminant
- This mathematical construction seems a very natural one. From a physical point of view, it makes contact with various condensed-matter systems: supersymmetric t - J [Wiegmann:1988, Sarkar:1991, Mavromatos:1999] quantum Hall effect [Hasebe], supersymmetric valence-bond solid states [Hasebe-Totsuka].

Superqubits and nonlocality

- Bell inequality

$$S(LHVM) \leq 2, \quad S(EPR) = 2\sqrt{2}$$

Superqubits and nonlocality

- Bell inequality

$$S(LHVM) \leq 2, \quad S(EPR) = 2\sqrt{2}$$

- It was later realised by Tsirelson [1980] that not only does the EPR-pair violate the Bell inequality, it violates it maximally: there is no quantum system that can do better.
- Is the Tsirelson Bound a fundamental limit of Nature? Can there at least in principle be a theory that crosses Tsirelson?

Superqubits and nonlocality

- Bell inequality

$$S(LHVM) \leq 2, \quad S(EPR) = 2\sqrt{2}$$

- It was later realised by Tsirelson [1980] that not only does the EPR-pair violate the Bell inequality, it violates it maximally: there is no quantum system that can do better.
- Is the Tsirelson Bound a fundamental limit of Nature? Can there at least in principle be a theory that crosses Tsirelson?
- In principle Nature could do better

$$S(\text{no-sulu}) \leq 4$$

without violating the no superluminal signaling principle [Popescu-Rohrlich:1994]

Superqubits and nonlocality

Can superqubits break Tsirlson's bound? [LB-Brádler-Duff:2012]

- Super resource:

$$\Gamma_{AB} = \left(1 + \frac{1}{2}X + \frac{3}{8}X^2\right) \left(\frac{1}{\sqrt{2}}(|00\rangle + |11\rangle) + \frac{p}{\sqrt{2}}\theta_A|\bullet 1\rangle + \frac{q}{\sqrt{2}}\theta_B|1\bullet\rangle\right)$$

- Local superunitaries:

$$Z_{iA} \otimes Z_{jB} = S(2r_i\theta_A)U(\alpha_i, \beta_i) \otimes S(2s_j\theta_B)U(\gamma_j, \delta_j).$$

- Born rule:

$$p_G(\varphi, \psi) = \langle \varphi | \psi \rangle (\langle \varphi | \psi \rangle)^\#.$$

The rationale behind this definition is clear: for ordinary (non-Grassmann) states we recover the usual Born rule.

- Grassmann norm:

$$|\tau|_R \stackrel{\text{df}}{=} \int \prod_i^n e^{\theta_i \theta_i^\#} \tau d^{2n}\theta,$$

Beats Tsirelson's bound - but at a cost: extended probabilities!

Super SLOCC algebras

Table : The action of the $\mathfrak{osp}(1|2)$ generators on the superqubit fields.

Generator	Field acted upon	
	a_{A_3}	a_{\bullet}
$P_{A_1 A_2}$	$\varepsilon_{(A_1 A_3} a_{ A_2)}$	0
$2Q_{A_1}$	$\varepsilon_{A_1 A_3} a_{\bullet}$	a_{A_1}

Table : The action of the $\mathfrak{osp}(1|2) \oplus \mathfrak{osp}(1|2)$ generators on the 2-superqubit fields.

Generator	Field acted upon			
	Bosons		Fermions	
	$a_{A_3 B_3}$	$a_{\bullet\bullet}$	$a_{A_3\bullet}$	$a_{\bullet B_3}$
$P_{A_1 A_2}$	$\varepsilon_{(A_1 A_3} a_{ A_2) B_3}$	0	$\varepsilon_{(A_1 A_3} a_{ A_2)\bullet}$	0
$P_{B_1 B_2}$	$\varepsilon_{(B_1 B_3} a_{A_3 B_2)}$	0	0	$\varepsilon_{(B_1 B_3} a_{\bullet B_2)}$
$2Q_{A_1}$	$\varepsilon_{A_1 A_3} a_{\bullet B_3}$	$a_{A_1\bullet}$	$\varepsilon_{A_1 A_3} a_{\bullet\bullet}$	$a_{A_1 B_3}$
$2Q_{B_1}$	$\varepsilon_{B_1 B_3} a_{A_3\bullet}$	$-a_{\bullet B_1}$	$a_{A_3 B_1}$	$-\varepsilon_{B_1 B_3} a_{\bullet\bullet}$

Super SLOCC algebras

Table : The action of the $\mathfrak{osp}(1|2) \oplus \mathfrak{osp}(1|2) \oplus \mathfrak{osp}(1|2)$ generators on the 3-superqubit fields.

Generator	Bosons acted upon			
	$a_{A_3} B_3 C_3$	$a_{A_3} \bullet \bullet$	$a_{\bullet} B_3 \bullet$	$a_{\bullet \bullet} C_3$
$P_{A_1 A_2}$	$\varepsilon(A_1 A_3 a A_2) B_3 C_3$	$\varepsilon(A_1 A_3 a A_2) \bullet \bullet$	0	0
$P_{B_1 B_2}$	$\varepsilon(B_1 B_3 a_{A_3} B_3) C_2$	0	$\varepsilon(B_1 B_3 a_{\bullet} A_2) \bullet$	0
$P_{C_1 C_2}$	$\varepsilon(C_1 C_3 a_{A_3} B_3 C_2)$	0	0	$\varepsilon(C_1 C_3 a_{\bullet \bullet} C_2)$
$2Q_{A_1}$	$\varepsilon_{A_1 A_3} a_{\bullet} B_3 C_3$	$\varepsilon_{A_1 A_3} a_{\bullet \bullet \bullet}$	$a_{A_1} B_3 \bullet$	$a_{A_1} \bullet C_3$
$2Q_{B_1}$	$\varepsilon_{B_1 B_3} a_{A_3} \bullet C_3$	$a_{A_3} B_1 \bullet$	$-\varepsilon_{B_1 B_3} a_{\bullet \bullet \bullet}$	$-a_{\bullet} B_1 C_3$
$2Q_{C_1}$	$\varepsilon_{C_1 C_3} a_{A_3} B_3 \bullet$	$-a_{A_3} \bullet C_1$	$-a_{\bullet} B_3 C_1$	$\varepsilon_{C_1 C_3} a_{\bullet \bullet \bullet}$
	Fermions acted upon			
	$a_{A_3} B_3 \bullet$	$a_{A_3} \bullet C_3$	$a_{\bullet} B_3 C_3$	$a_{\bullet \bullet \bullet}$
$P_{A_1 A_2}$	$\varepsilon(A_1 A_3 a A_2) B_3 \bullet$	$\varepsilon(A_1 A_3 a A_2) \bullet C_3$	0	0
$P_{B_1 B_2}$	$\varepsilon(B_1 B_3 a_{A_3} B_3) \bullet$	0	$\varepsilon(B_1 B_3 a_{\bullet} B_3) C_2$	0
$P_{C_1 C_2}$	0	$\varepsilon(C_1 C_3 a_{A_3} \bullet C_2)$	$\varepsilon(C_1 C_3 a_{\bullet} B_3 C_2)$	0
$2Q_{A_1}$	$\varepsilon_{A_1 A_3} a_{\bullet} B_3 \bullet$	$\varepsilon_{A_1 A_3} a_{\bullet \bullet \bullet} C_3$	$a_{A_1} B_3 C_3$	$a_{A_1} \bullet \bullet$
$2Q_{B_1}$	$\varepsilon_{B_1 B_3} a_{A_3} \bullet \bullet$	$a_{A_3} B_1 C_3$	$-\varepsilon_{B_1 B_3} a_{\bullet \bullet \bullet} C_3$	$-a_{\bullet} B_1 \bullet$
$2Q_{C_1}$	$a_{A_3} B_3 C_1$	$-\varepsilon_{C_1 C_3} a_{A_3} \bullet \bullet$	$-\varepsilon_{C_1 C_3} a_{\bullet} B_3 \bullet$	$a_{\bullet \bullet} C_1$